

# CHAPTER 9

## GRAPHS

### 9.1 A GENTLE INTRODUCTION

# FOUR PROBLEMS



The Bridges of Königsberg



Three House-Three Utility



Four Color

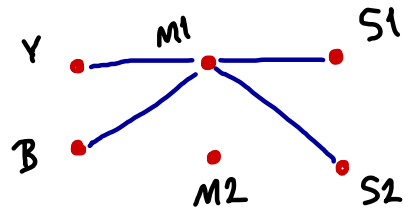


Traveling Salesman

# A SAMPLE PROBLEM

Among you, your buddy, two mothers, and two sisters, some people hug. There are no hugs between buddies, mothers, or sisters. The other 5 people tell you they all hugged different numbers of people. How many people did you hug?

Your buddy did not hug 4 people (otherwise no one hugged 0 people)  $\leadsto$  may assume your mother hugged 4



So your buddy's mom hugged 0.  
If your buddy hugged 3, nobody hugged 1.  
If your buddy hugged 1, nobody hugged 3.  
 $\leadsto$  your buddy hugged 2  $\leadsto$  you hugged 2.

## 9.2 DEFINITIONS AND BASIC PROPERTIES

# GRAPHS

A **graph** is a pair of sets  $V$  and  $E$ , where  $V \neq \emptyset$  and each element of  $E$  is a pair of elements of  $V$ .

Write  $G = G(V, E)$ .

For us, graphs are **finite**, that is,  $|V|$  is finite.

The elements of  $V$  and  $E$  are called **vertices** and **edges**.

**EXAMPLE.**  $V =$  Facebook users  
 $E =$  Friendships

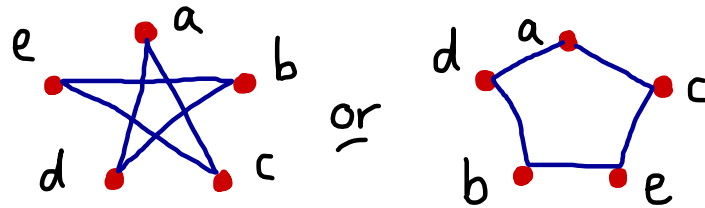
# GRAPHS

We can represent graphs with pictures.

**EXAMPLE.** Consider the graph  $G(V, E)$  where

$$V = \{a, b, c, d, e\}$$

$$E = \{\{d, b\}, \{a, c\}, \{e, b\}, \{e, c\}, \{d, a\}\}$$



Can describe a graph with a picture instead of set notation.

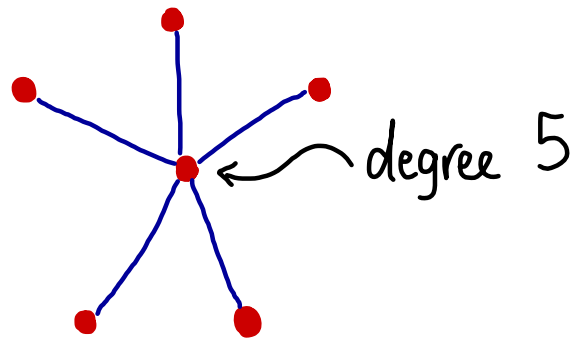
Could also write  $E = \{db, ac, eb, ec, da\}$ .

We say  $a$  is adjacent to  $c$  and  $d$ , and  $ac$  is incident to  $a$  and  $c$ .

# DEGREES

The **degree** of a vertex  $v$  is the number of edges incident to  $v$ . Write **deg  $v$** .

If  $\text{deg } v = 0$ , we say  $v$  is **isolated**.



# PSEUDOGRAPHS

The following two phenomena are not allowed in a graph:



If we allow these, we get what is called a **pseudograph**.


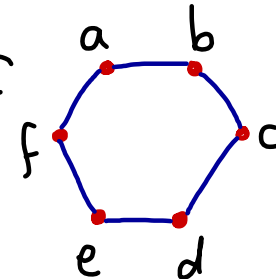
Pseudographs are harder to write down with set notation, so we usually describe them with a picture.

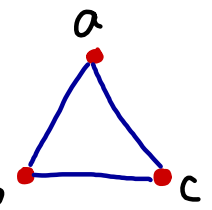
**EXAMPLE.** Vertices are web pages  
Edges are links



# SUBGRAPHS

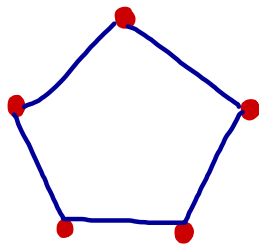
A **subgraph** of a graph  $G(V, E)$  is a graph  $G(V', E')$  where  
 $V' \subseteq V$  and  
 $E' \subseteq E$

EXAMPLE.  is a subgraph of 

but  is not.

Also: Can delete any number of edges to get a subgraph.  
Can delete any number of vertices (and all incident edges) to get a subgraph.

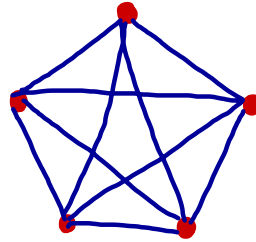
# THREE SPECIAL FAMILIES



$C_5$

$C_n$

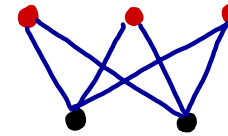
$n$ -cycle



$K_5$

$K_n$

complete graph



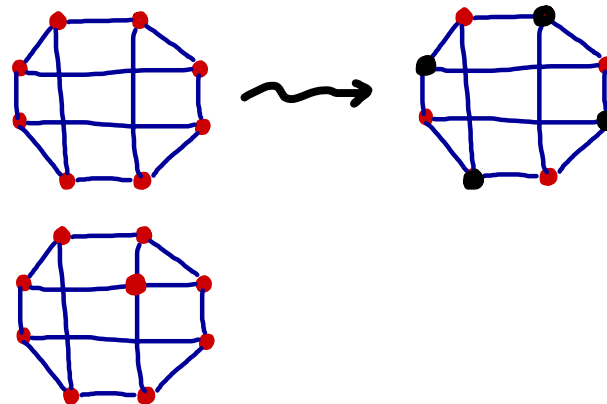
$K_{3,2}$

$K_{m,n}$

Complete bipartite  
graph

# BIPARTITE GRAPHS

A **bipartite** graph is one whose vertex set can be partitioned into two sets  $V_1$  and  $V_2$  so that each edge joins an element of  $V_1$  to an element of  $V_2$ .



**FACT.** A bipartite graph contains no triangles. More generally, a bipartite graph contains no odd cycles.

# THE HANDSHAKING LEMMA

**PROPOSITION.** The sum of the degrees of the vertices of a pseudograph is an even number.  
Specifically:

$$\sum_{v \in V} \deg v = 2|E|$$



Leonhard Euler

**HANDSHAKING LEMMA.** The number of odd degree vertices of a pseudograph is even.

**PROOF.** 
$$\sum_{v \in V} \deg v = \sum_{v \text{ even}} \deg v + \sum_{v \text{ odd}} \deg v$$

Revisit the hugging problem.

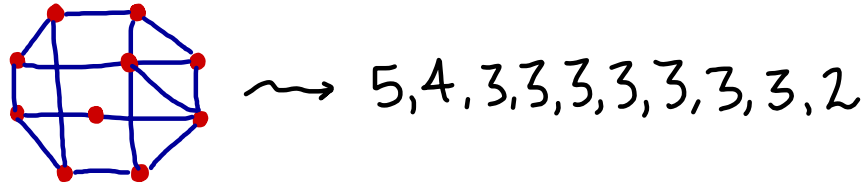
# THE HANDSHAKING LEMMA

**PROBLEM.** A graph has 50 edges, 4 vertices of degree 2, 6 of degree 5, 8 of degree 4, all other vertices have degree 6. How many vertices does the graph have?

**PROBLEM.** Out of 24 curling players, 78 pairs have played on the same team. Show that one has played on the same team as 7 others. Show that one has played on the same team with no more than 6 others.

# DEGREE SEQUENCE

Say  $d_1, \dots, d_n$  are the degrees of the vertices of a pseudograph, where  $d_1 \geq d_2 \geq \dots \geq d_n$ . Then  $d_1, \dots, d_n$  is the **degree sequence** of the pseudograph.



## 9.3 GRAPH ISOMORPHISM

# GRAPH ISOMORPHISM

Two graphs  $G(V,E)$  and  $G(V',E')$  are *isomorphic* if there is a bijection

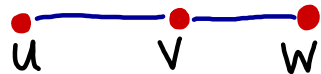
$$V \rightarrow V'$$

that preserves adjacency and nonadjacency.

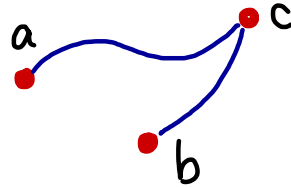
In other words, two graphs are isomorphic if there is a change of labels taking one to the other.

**EXAMPLE.**

$$V = \{u, v, w\}$$
$$E = \{uv, vw\}$$



$$V' = \{a, b, c\}$$
$$E' = \{ac, cb\}$$



$$V \rightarrow V'$$

$$u \mapsto a$$

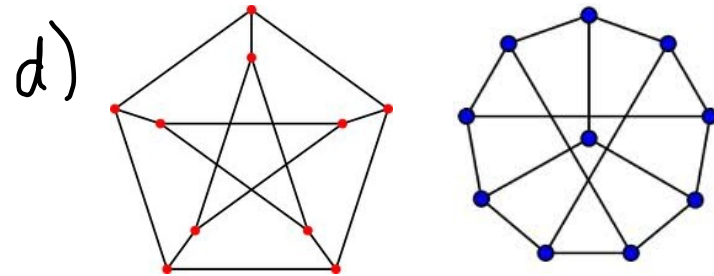
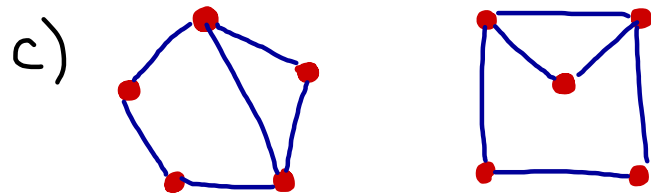
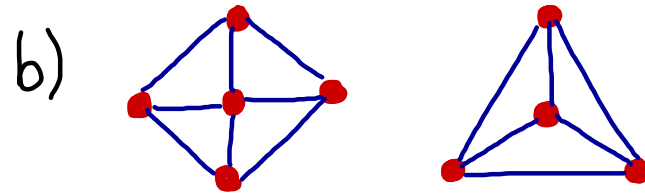
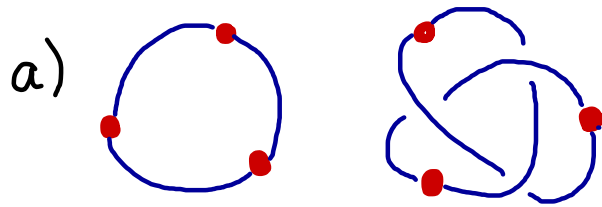
$$w \mapsto b$$

$$v \mapsto c$$



# GRAPH ISOMORPHISM

Which of the following pairs are isomorphic?

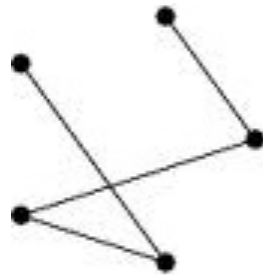


# INVARIANTS OF GRAPHS

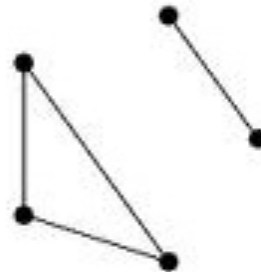
We can use the following "fingerprints" of graphs in order to tell if two graphs are *different*:

- (i) Number of vertices
- (ii) Number of edges
- (iii) Degree sequence
- etc.

It is possible for two graphs to have the same degree sequence and be nonisomorphic:



{2, 2, 2, 1, 1}



{2, 2, 2, 1, 1}

# EXAMPLES

Which of the following graphs are isomorphic?

