

# JOHNSON I

Thm  $I(S_g)$  is fin. gen. by Dehn twists for  $g \geq 3$ .

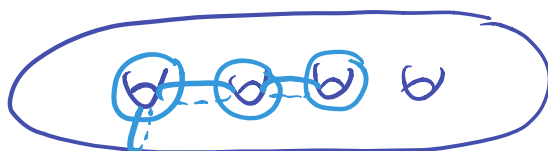
**Basic strategy**

1. List prospective generators  $\{g_i\}$ .  
s.t.  $g_i$  is a BP map of genus 1.
2. Show  $\langle g_i \rangle \trianglelefteq \text{Mod}(S_g)$ .

This suffices since  $\langle\langle g_i \rangle\rangle_{\text{Mod}(S_g)} = I(S_g)$ .

Chains and BP maps.

A chain:



$\rightsquigarrow$  BP map.

Given a chain, can resolve intersections to get another chain. Can also take subchains.

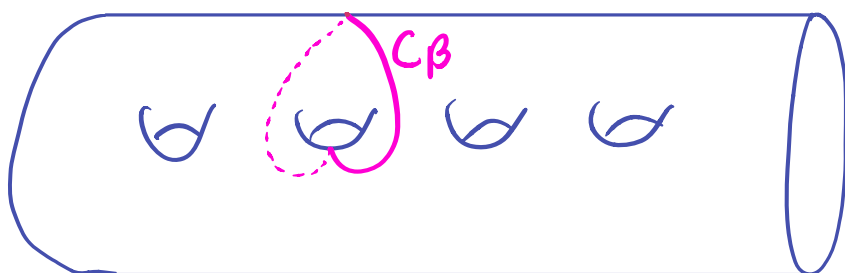
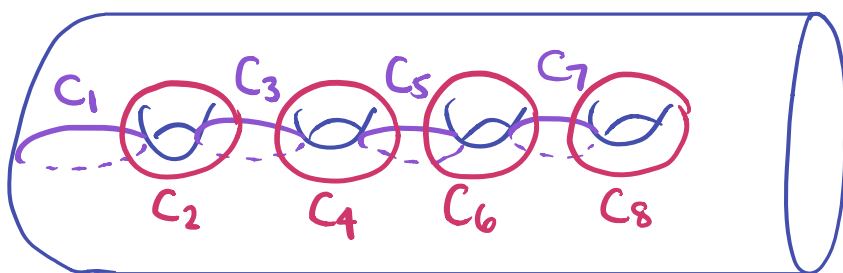
Given a chain  $ch(c_1, \dots, c_n)$

$ch(i_1, i_2, \dots, i_{k+1})$

denotes the chain you get by combining  $c_{i_1}, \dots, c_{i_2-1}$   
 $c_{i_2}, \dots, c_{i_3-1}$  etc. dropping  $c_{k+1}, \dots, c_n$ . "subchain"

Denote the BP-map  $[i_1, i_2, \dots, i_{k+1}]$

# LISTING THE GENERATORS



Consider the chains:

$(C_1, \dots, C_{2g})$  straight chain

$(C_\beta, C_5, \dots, C_{2g})$   $\beta$ -chain

Use same notation for subchains of  $\beta$ -chain:

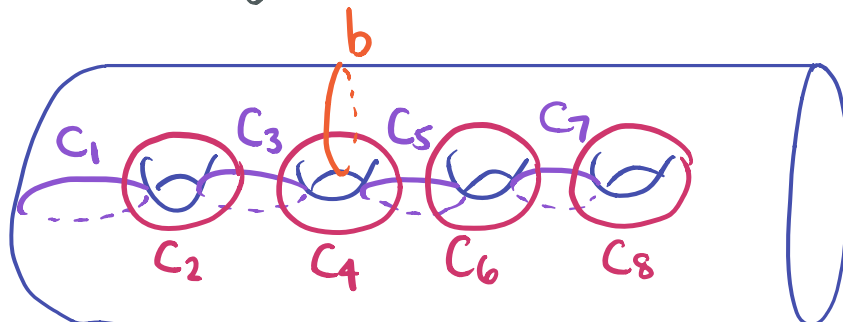
$(\beta_i) = \text{surger } C_\beta, C_5, \dots, C_{i-1}$

Theorem. For  $g \geq 3$  the odd subchain maps of straight chain &  $\beta$ -chain generate  $I(S'_g)$

Since  $I(S'_g) \rightarrow I(S_g)$  this gives closed case as well.

SETUP.

Humphries generators:



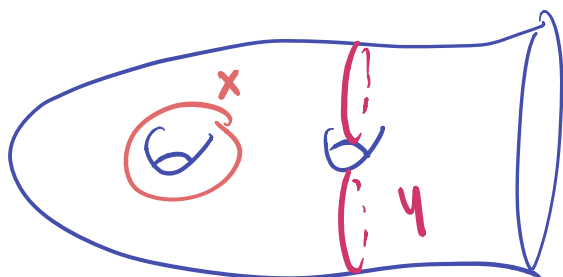
Let  $J(S'_g)$  &  $J(S_g)$  denote groups gen by Johnson's generators.

As above, need to show

$$T_x * \gamma = T_x \gamma T_x^{-1} \in J(S_g)$$

$\forall x \in$  Humphries set  
 $\gamma \in$  Johnson set

In many cases  $T_x * \gamma$  equals  $\gamma$  or is another Johnson gen:



# CONJUGATING BY A POSITIVE TWIST IN THE CHAIN

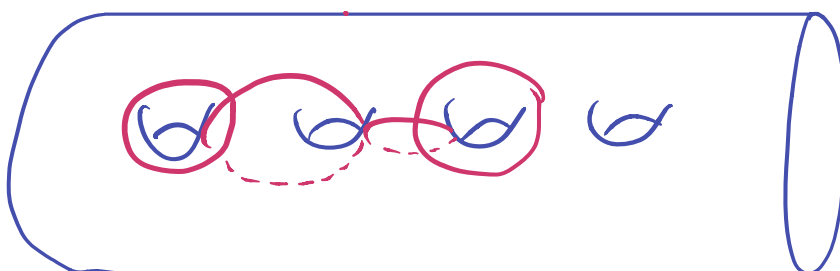
Prop.

$j$ in $\{i_1, \dots\}$	$j+1$ in $\{i_1, \dots\}$	$T_{C_j} * [i_1 i_2 \dots]$
✓	✓	$[i_1 i_2 \dots]$
✗	✗	$[i_1 i_2 \dots]$
$i_m$	✗	$[i_1 \dots i_{m-1} i_{m+1} \dots]$
✗	$i_m$	$[i_1 \dots i_{m-1} i_{m-1} \dots]$

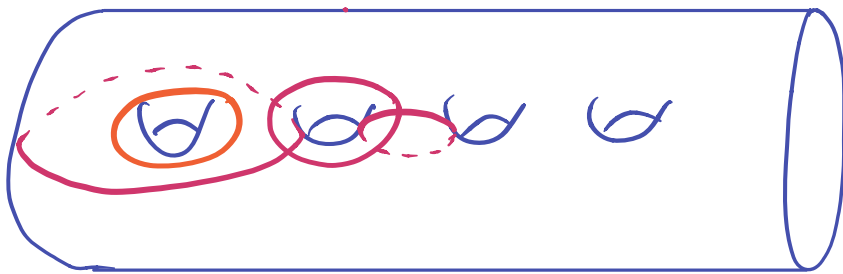
The lemma completely characterizes commuting among straight Johnson & Humphries gens.

PF (by examples)

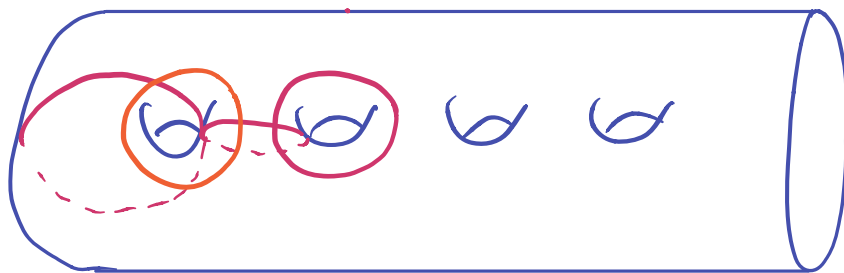
Example  $j=2$   $[2 3 5 6]$



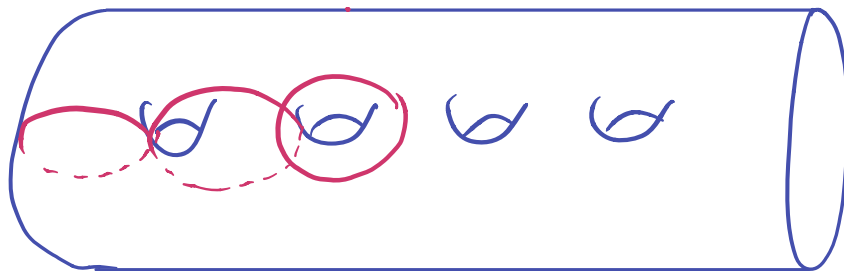
Example  $j=2$  [1 4 5 6]



Example  $j=2$  [1 3 4 5]



[1 2 4 5]



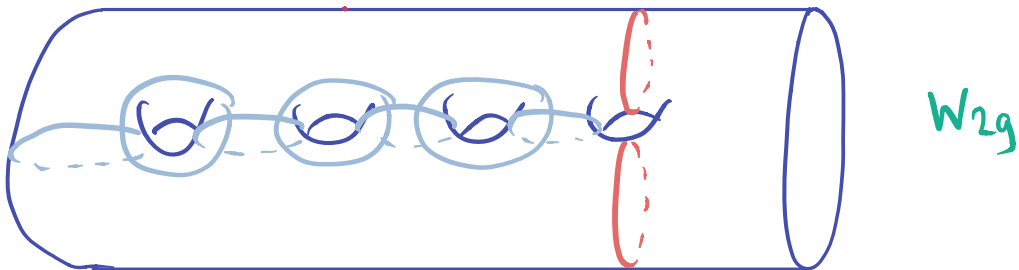
Third case similar.

Upshot of Lemma: If we conj a chain map by a positive twist about a curve in the chain we get a subchain map. What about negative twists?

# GENERATING THE KERNEL OF $I(S'_g) \rightarrow I(S_g)$

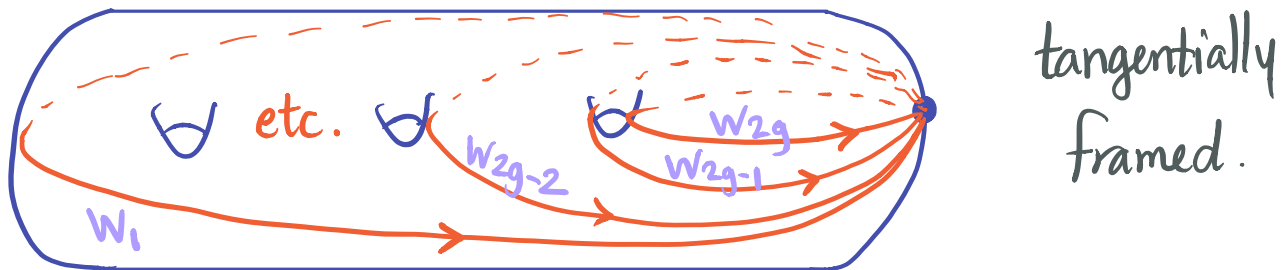
Define  $W_i = [1 \ 2 \ 3 \ \dots \ \hat{i} \ \dots \ 2g+1]$   
 for  $i \in \{1, \dots, 2g+1\}$ . "maximal odd subchain maps"

Each lies in  $\pi_1 \text{UT}(S_g)$



Lemma.  $\pi_1 \text{UT}(S_g)$  is gen by  $W_1, \dots, W_{2g}, T_b * W_1$

Pf. The corresponding push maps are



These clearly generate  $\tilde{\pi}_1$ . Need to get the fiber of  $\tilde{\pi}_1 \text{UT} \rightarrow \pi_1$ .

Claim.  $T_b * W_1 = W_4 W_3^{-1} W_2 T_b$  □

Cor.  $\pi_1 \text{UT}(S'_g) \leq J(S'_g)$

If. Use lantern to get  $T_b * W_1$  (Cor 1 in paper).

# CONJUGATING BY A NEGATIVE TWIST IN THE CHAIN

PROP. Fix a chain  $(c_i)$ . The subgroup of  $I(S_g')$   
 $\langle \text{odd subchain maps of length } 2k-1 \rangle$   
is normalized by  $\langle T_{c_i} \rangle$ .

Lemma.  $\langle W_1, \dots, W_{2g+1} \rangle$  is normalized by  
 $\langle T_{c_1}, \dots, T_{c_{2g}} \rangle$

Pf. Need to check  $T_{c_i}^{\pm 1}(W_j) \in \langle W_1, \dots, W_{2g+1} \rangle$   
in  $\Pi_1 UT(S_g')$ .

Pf of Prop. Let  $f = [i_1 i_2 \dots i_{2k}]$

Fix a  $c_j$  in original chain.

As above  $T_{c_j} * f = f$  unless

exactly one of  $j, j+1$  is an  $i_\square$ .

Enlarge  $(i_1, i_2, \dots, i_{2k})$  to include both  $j$  &  $j+1$ .

$\rightsquigarrow$  subchain of length  $2k$   $(c'_i)$  w/  $c_j$  basic

Lemma  $\Rightarrow T_{c_j}^{\pm 1} * f$  is a product of maximal  
subchain maps of  $(c'_i)$ . But these  
are all subchains of original chain  $\square$

## TWO MORE TOOLS.

### ① INDUCTION

If our Johnson & Humphries gens both lie in  $S_{g-1}'$ , can apply induction. Works because both sets of gens agree with  $S_{g-1}' \hookrightarrow S_g'$  (this is one specific inclusion).

N.B. Base case is subgp of  $I(S_2')$  gen. by BP maps i.e.  $\pi_1 \text{UT}(S_2)$   
i.e.  $\langle W_1, \dots, W_5 \rangle$ .

### ② CONTROLLED CHANGE OF COORDS

Suppose  $T_{c_i} \in \text{Humphries}$   
 $f \in \text{Mod}(S_g')$  already known to  
normalize  $J(S_g')$ .

$h \in \text{Johnson}$

Then checking  $T_{c_i} * h \in J(S_g')$  is equiv. to  
checking  $f * (T_{c_i} * h)$   
 $= T_{c_i} * (f * h) \in J(S_g')$



# PROOF OF THE THEOREM

Need to conjugate all Johnson gens by  $T_{c_i}^{\pm 1}, T_b^{\pm 1} \in \text{Humphries}$ .

Step 1.  $T_{c_1}, T_{c_2}, T_{c_5}, \dots, T_{c_{2g}}$  normalize  $J(S'_g)$ .

$T_{c_1}, T_{c_2}$  in the straight chain, disjoint from  $\beta$ -chain.

$T_{c_5}, \dots, T_{c_{2g}}$  in both chains.

Apply Prop.

Step 2.  $T_{c_3}$  normalizes  $J(S'_g)$ .

$T_{c_3}$  normalizes straight chains by Prop.

By controlled change of coords, can restrict to consecutive  $\beta$ -chain maps.

$[\beta \ 5 \ 6 \ \dots]$ .

If chain not maximal, use induction.

If it is, need a relation (Cor 2 in paper).

But this is one check.

Step 3.  $T_{c_4}$  normalizes  $J(S'_g)$  similar to  $T_{c_3}$

Step 4.  $T_b$  normalizes  $J(S'_g)$ .

$\beta$ -chains: controlled change of coords  
reduces consecutive  $\beta$ -chains &  
induction reduces to

$$T_b^{\pm 1} * [\beta \ 5 \ \dots \ 2g+1] \quad \text{single check}$$

Straight chains  $[i_1 \ \dots \ i_{2k}]$

Can assume  $i_1 \leq 4$   $i_{2k} \geq 5$   
(otherwise  $T_b$  commutes)

Change of coords reduces to: consecutive  
straight chain starting with 1, 2, 3, or 4.  
and going up to  $2g$  or  $2g+1$  (induction).

If it starts with 1 or 2 it is maximal,  
i.e. in  $\Pi_1 U T$ .

Only two remaining cases:

$$[4 \ 5 \ \dots \ 2g+1] \ \& \ [3 \ 4 \ \dots \ 2g].$$

