

## INGREDIENTS FOR ACYLINDRICITY

Thm 1.  $d(a, b) \geq 3 \Rightarrow |\text{Stab}_{\text{MCG}}(a) \cap \text{Stab}_{\text{MCG}}(b)| \leq N_0 = N_0(S)$

Pf idea.  $a \cup b \rightsquigarrow \text{cell decomp of } S$

topological lemma: any  $f \in \text{Stab}(a) \cap \text{Stab}(b)$  has a rep that preserves the cell decomp.

- The resulting auto. of the cell decomp is determined by where it sends one 2-cell.

But the number of nonrectangular 2-cells is at least one and is bounded by a fn of  $S$ .  $\square$

  $G(a, b; r) = \text{curves that lie on some tight geod. from } a' \text{ to } b'$   
where  $d(a, a') \leq r, d(b, b') \leq r$ .

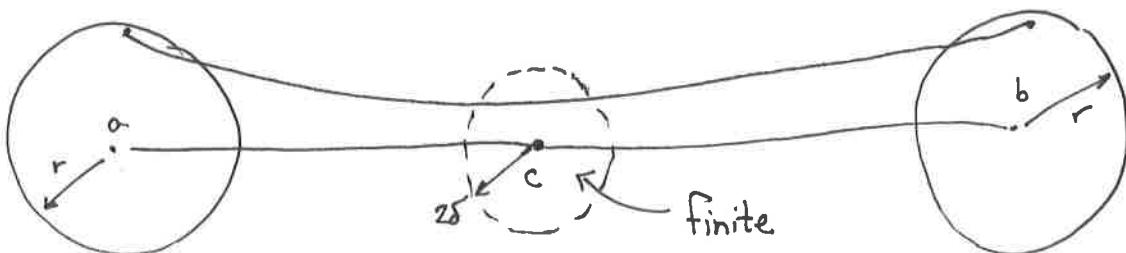
Thm 2. Fix  $r \geq 0$ .

$a, b \in C(S)$  with  $d(a, b) \geq 2r + 2(10\delta + 1) + 1$

$c \in \pi = \text{geod. from } a \text{ to } b$ .

$c \notin N_{r+10\delta+1}(a) \cup N_{r+10\delta+1}(b)$

$$\rightsquigarrow |G(a, b; r) \cap N_{2\delta}(c)| \leq D = D(S)$$



## Proof of ACYLINDRICITY

$$R = 4r + 24\delta + 7$$

$$N = N_0 (2r + 4\delta + 1)(8\delta + 7)D$$

Say  $d(a, b) \geq R$

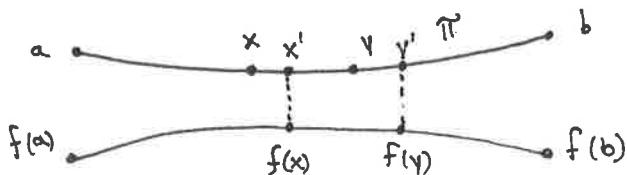
Pick  $x, y \in \pi = \text{tight geod from } a \text{ to } b$ .

s.t. ①  $d(x, y) = 3$

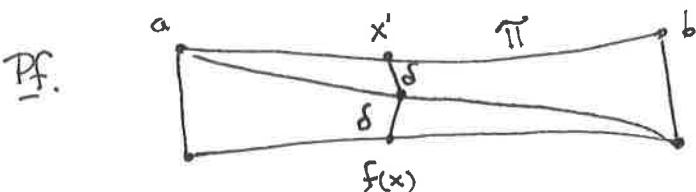
②  $d(\{x, y\}, \{a, b\}) \geq r + (10\delta + 1) + (2\delta + r) + 1$

Say  $f \in \text{MCG}(S)$  with  $d(a, f(a)) \leq r$ ,  $d(b, f(b)) \leq r$

Let  $x', y'$  proj's of  $\overset{f(x), f(y)}{\cancel{xy}}$  to  $\pi$ .



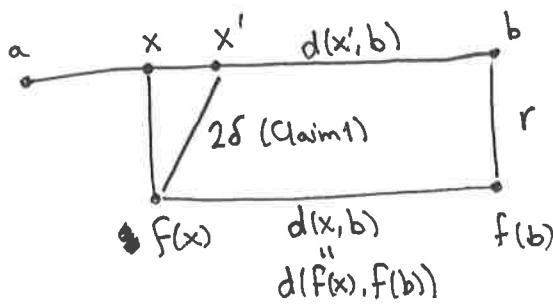
Claim 1.  $d(f(x), \pi) \leq 2\delta$ ,  $d(f(y), \pi) \leq 2\delta$



Use  $\delta$ -thinness plus fact that  $f(x)$  is far from the vertical sides.

Claim 2.  $d(x, x') \leq r + 2\delta$   $d(y, y') \leq r + 2\delta$

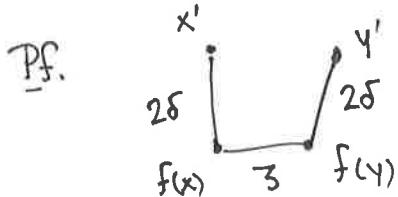
Pf. Assume  $x'$  to right of  $x$ :



$$\begin{aligned} d(x, x') &= d(x, b) - d(x', b) \\ &\leq (2\delta + d(x', b) + r) - d(x', b) \\ &= 2\delta + r \end{aligned}$$

If  $x'$  to left of  $x$ , replace  $b$  with  $a$ .

Claim 3.  $d(x', y') \leq 4\delta + 3$



Claim 4.  $d(x', a)$ ,  $d(y', b) \geq r + 10\delta + 2$

Pf. Immediate from Claim 2 & choice of  $x, y$ .

Claim 5. At most  $2r + 4\delta + 1$  choices for  $x'$ .

Pf. Immediate from Claim 2.

Claim 6. Given  $x'$ , at most:
 

- $(2r + 4\delta + 1)D$  choices for  $f(x)$ . Claim 4 + Thm 2
- $8\delta + 7$  choices for  $y'$  (Claim 3)
- $(8\delta + 7)D$  choices for  $f(y)$  Claim 4 + Thm 2

Acylindricity now follows from Thm 1, with  $N$  as above. □

## BOTTLENECKS

Remains to prove Thm 2. Here is a simpler version.

Thm.  $a, b \in C(S)$

$c \in \pi = \text{geod. from } a \text{ to } b.$

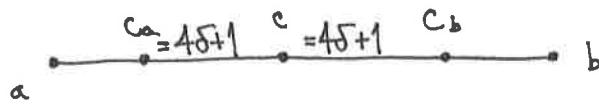
$$\leadsto |G(a, b) \cap N\delta(c)| \leq D.$$

$G(a, b) = G(a, b; 0) =$   
set of curves lying on  
some tight geod from  
 $a$  to  $b$ .

Pf. For simplicity, assume  $c$  is far from  $a, b$ :

$$d(c, \{a, b\}) \geq 4\delta + 1.$$

Choose  $c_a, c_b$ :



Enough to show that each elt of  $G(a, b) \cap N\delta(c)$  also lies on a tight filling multipath\* from  $c_a$  to  $c_b$  of length at most  $12\delta + 1$ .

Indeed, when we gave the algorithm for distance we showed there is a constant  $B = B(S, L)$  s.t. the number of curves that can lie on a tight filling multipath of length  $\leq L$  is bdd above by  $B$ .

\* A tight path  $(v_i)$  where  $|i-j| \geq 3 \Rightarrow v_i, v_j$  fill.